

# On the Difficulty of Nearest Neighbor Search

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## Introduction

### Large Scale NN Search in the Era of Big Data

- Big data: web multimedia, enterprise data centers, mobile/surveillance sensor systems, network nodes, etc....
- Large scale NN search for many big data applications
  - Retrieval from massive data such as multimedia search
  - Build neighborhood graphs for learning tasks like spectral clustering
  - ...

### Large Scale NN Search Methods

- Exhaustive NN search: prohibitively expensive for large scale data
- Recently many approximate NN search Methods
  - Tree based methods: kd-tree, metric tree, ...
  - Hashing based methods: Locality Sensitive Hashing (LSH), spectral hashing, ...

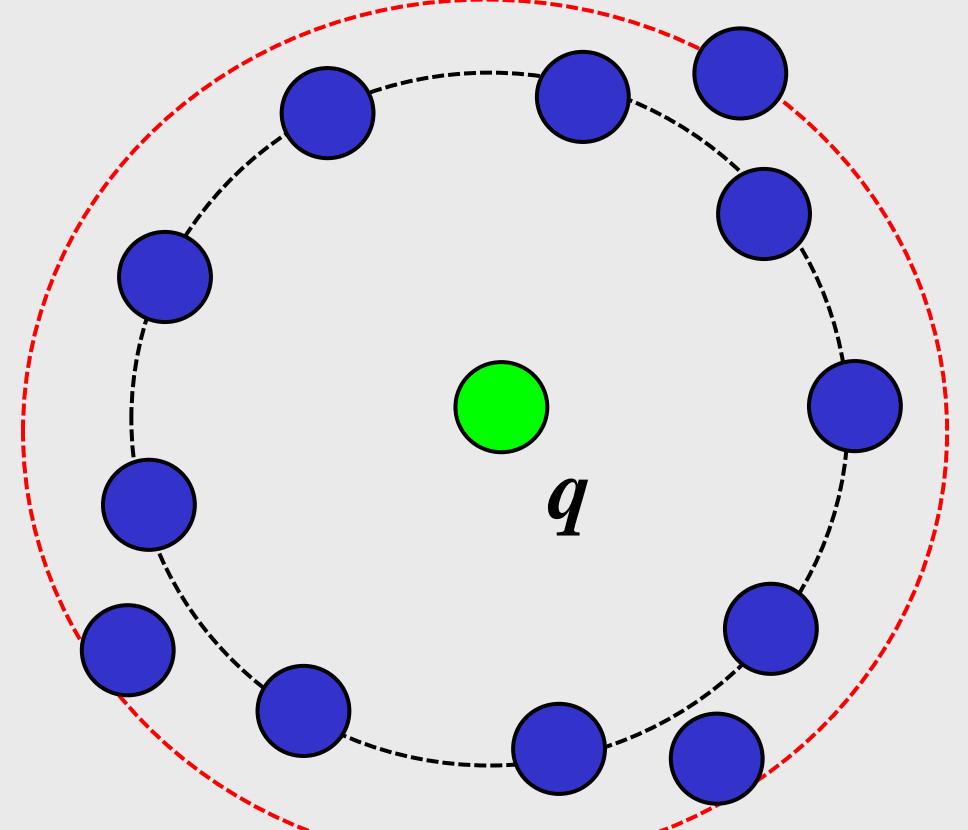
### A More Fundamental Problem

- How to measure the difficulty of a given data set for NN search, independent of NN search Methods?
- Moreover, what data properties affect the difficulty, and how?

## Difficulty Measure— Relative Contrast

### A Toy Example

If we can not differentiate NN point from other points, NN search is not meaningful!



### A Concrete Measure -- Relative Contrast

$$\text{Relative Contrast} = C_r = \frac{D_{\text{random}}}{D_{\text{nn}}} = \frac{E_x[D(q, x)]}{D(q, x_{\text{nn}})}$$

$$C_r = \frac{E_{q,x}[D(q, x)]}{E_q[D(q, x_{\text{nn}})]}$$

High Relative Contrast  $\rightarrow$  more meaningful search  
If  $C_r \rightarrow 1$ , search not meaningful

### Normalized Variance $\sigma'$

- Given a database  $X = \{x_i\}_{i=1}^n$ ,  $x \in \mathbb{R}^d$ , a query  $q$ , and a distance metric (say  $L_1$ ),

$$D(q, x) = \sum_{j=1}^d |q^j - x^j| \implies D = \sum_{j=1}^d D_j$$

- Let dimensions be i.i.d. with  $E[D_j] = \mu_d$ ,  $\text{var}[D_j] = \sigma_d^2$

From central limit theorem for large enough  $d$

$$D \sim N(\mu, \sigma^2) \quad \mu = d\mu_d \quad \sigma^2 = d\sigma_d^2$$

- If data is scaled such that  $\mu^2 = 1$ , then new variance

$$\sigma'^2 = \frac{\sigma^2}{\mu^2} = \frac{1}{d} \frac{\sigma_d^2}{\mu_d^2} \implies d \rightarrow \infty \implies \sigma'^2 \rightarrow 0$$

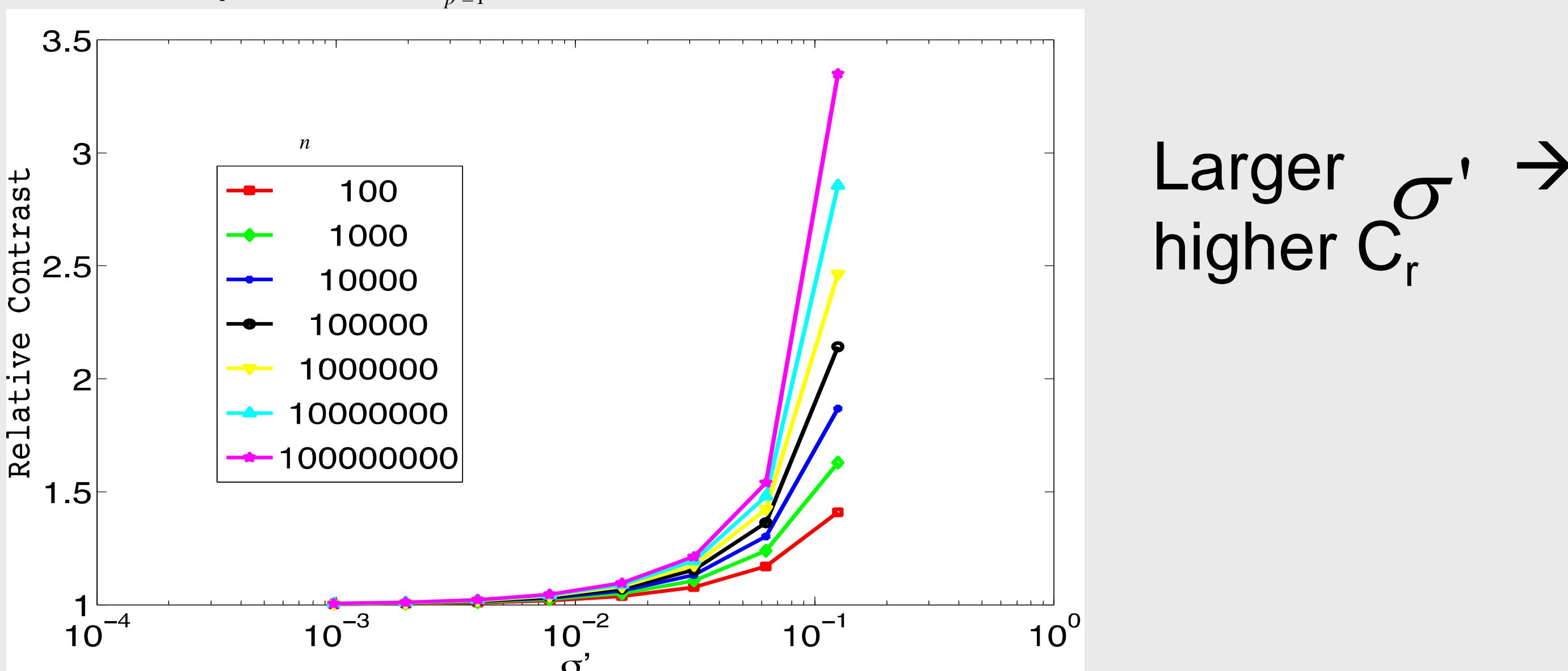
Distance to any point becomes roughly the same !

## What Affect NN Search and How?

### NN Search Difficulty — Relative Contrast $C_r$

$$C_r = \frac{D_{\text{mean}}}{D_{\text{min}}} \approx \frac{1}{[1 + \phi^{-1}(\frac{1}{n} + \phi(\frac{-1}{\sigma'}))\sigma']^{\frac{1}{p}}}$$

$\phi$  - standard Gaussian cdf



Larger  $n \rightarrow$  higher  $C_r$

### Normalized Variance $\sigma'$

Suppose dimensions are i.i.d., and each dimension has a probability  $s$  of being non-zero

Probability of both  $x^j$  and  $q^j$  being non-zero =  $s^2$

Probability of either  $x^j$  or  $q^j$  being non-zero =  $2s(1-s)$

For non-zero entries, let  $E[|x^j|^p] = m_p$ ,  $E[|q^j - x^j|^p] = m'_p$

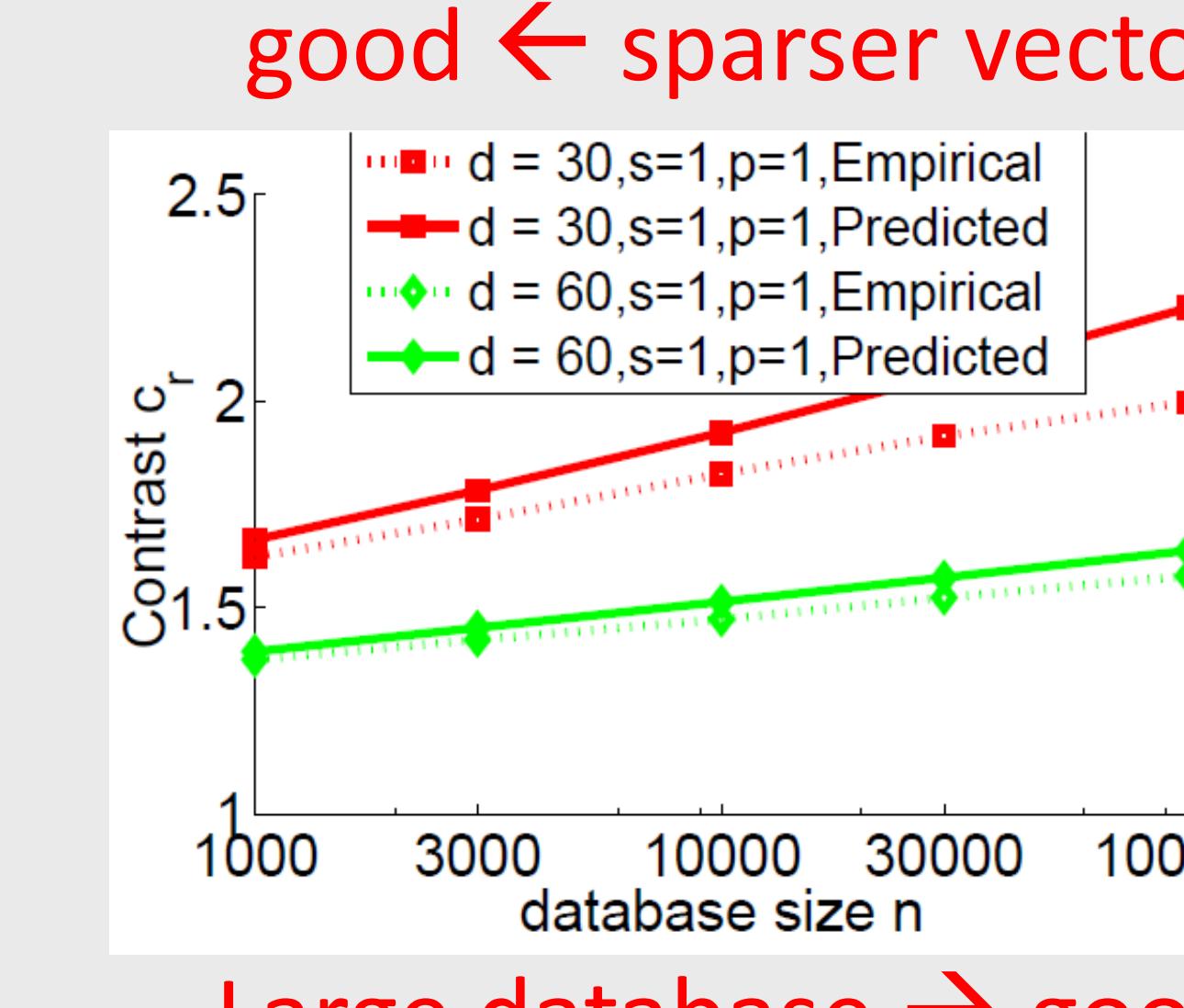
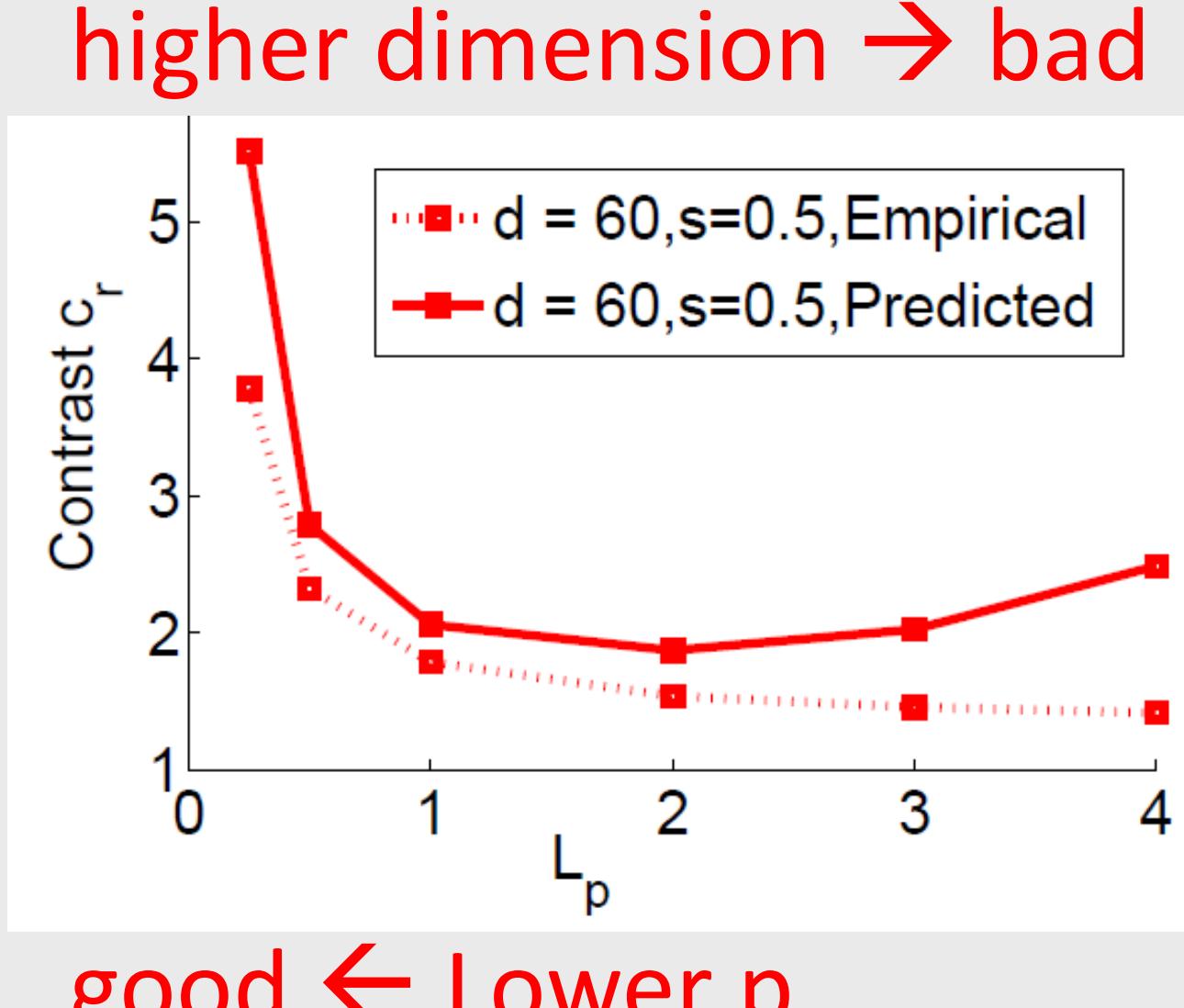
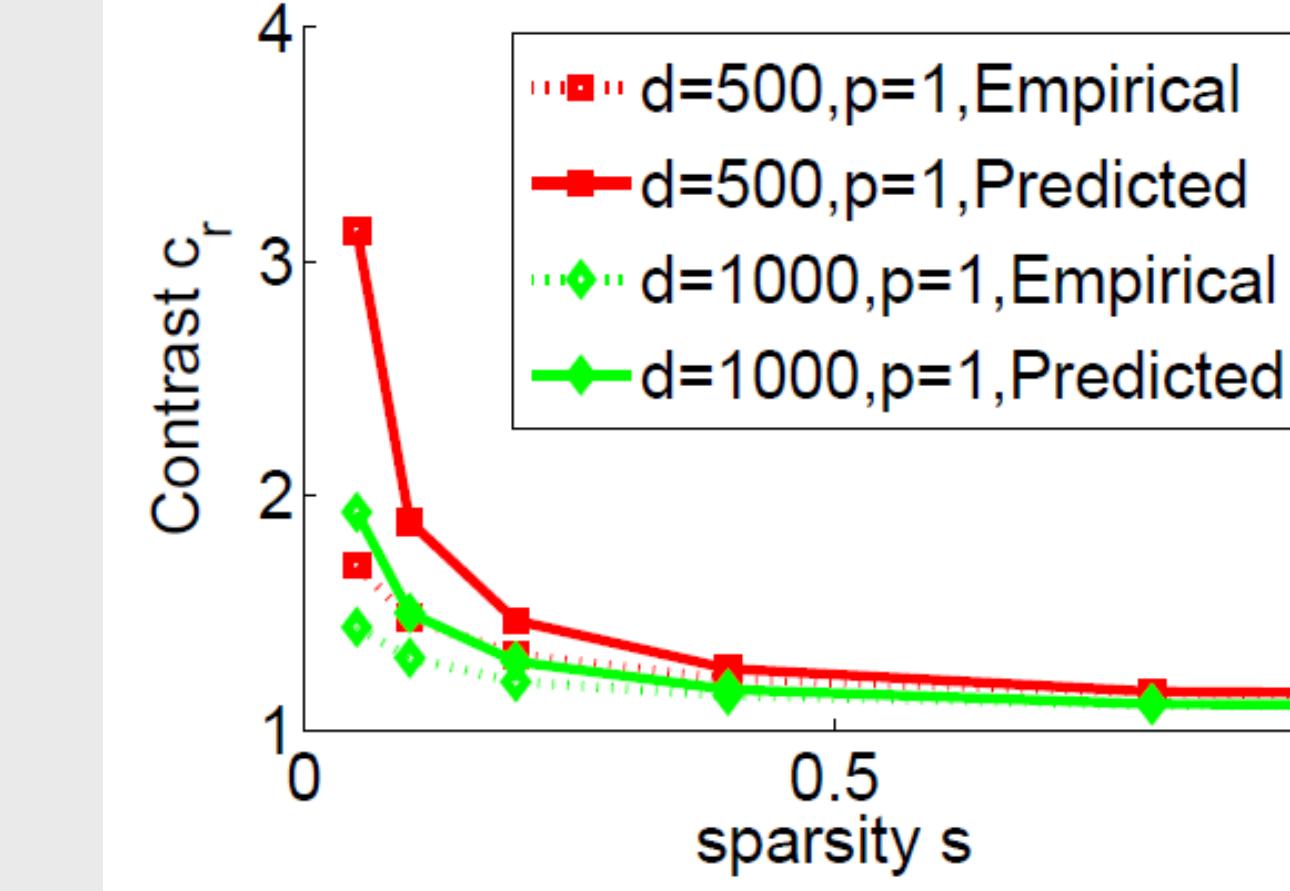
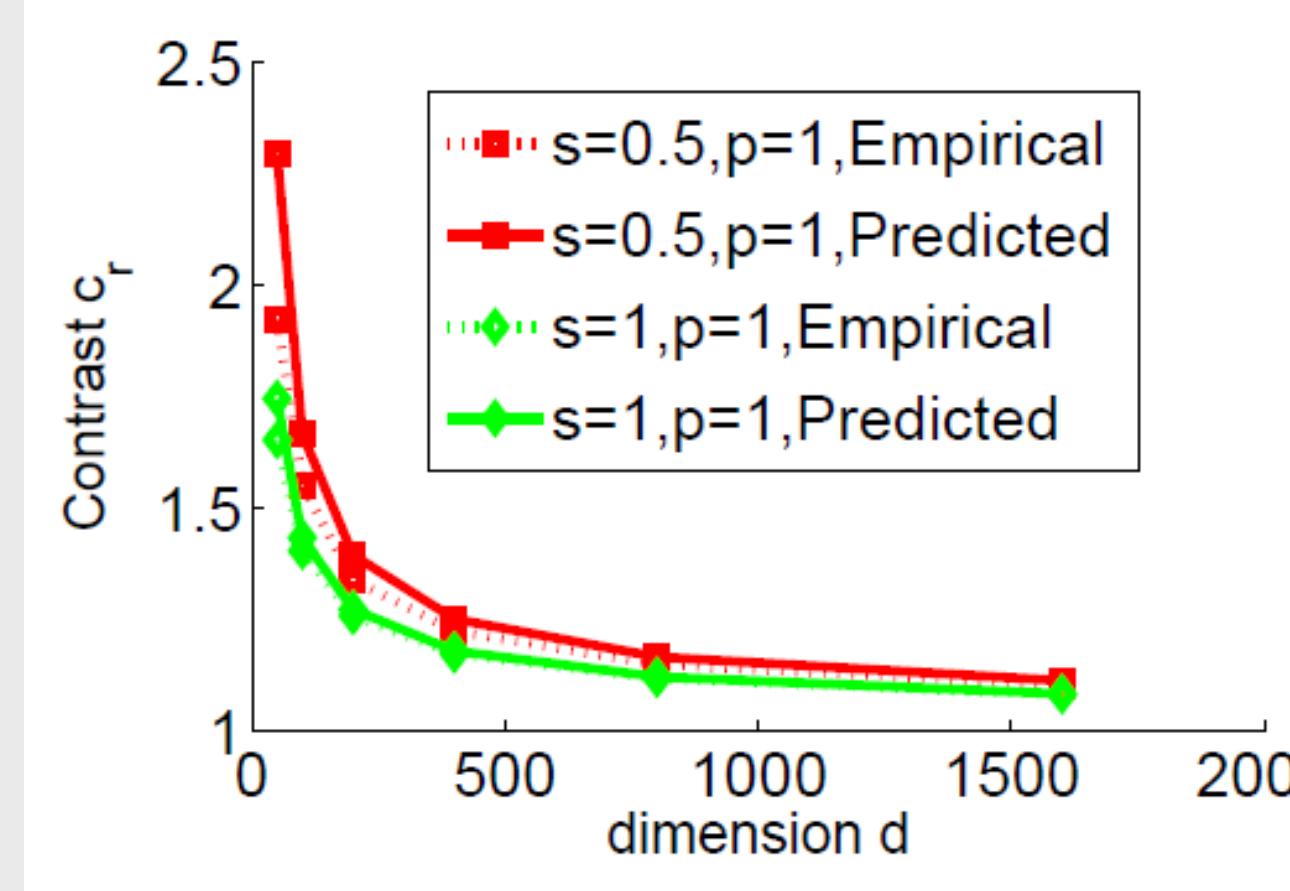
$$\mu_d = s^2 m_p + 2s(1-s)m_p \quad \sigma_d^2 = s^2 m_p + 2s(1-s)m_p - \mu_d^2$$

$$\sigma' = \frac{1}{d^{1/2}} \sqrt{\frac{s[(m'_p - 2m_p)s + 2m_p]}{s^2[(m'_p - 2m_p)s + 2m_p]^2}} - 1$$

### Data Properties:

dimension  $d$ , sparsity  $s$ , L<sub>p</sub> distance  $p$ , database size  $n$

### Experiments on Synthetic Data Sampled from U[0,1]



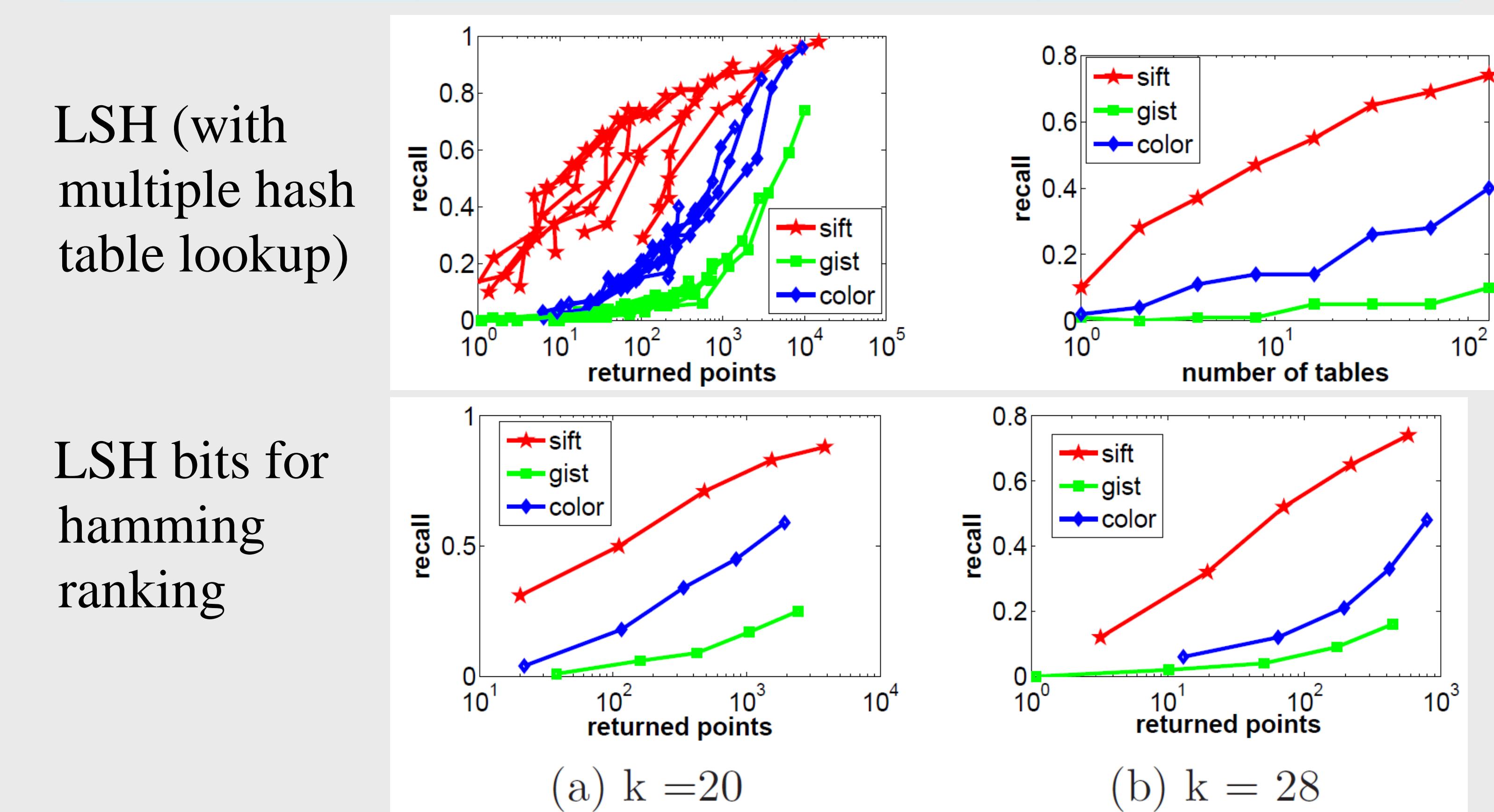
good ← Lower p      Large database → good

## Relative Contrast and LSH Complexity

**Theorem 3.1** LSH can find the exact nearest neighbor with probability  $1 - \delta$  by returning  $O(\log \frac{1}{\delta} n^{g(C_r)})$  candidate points, where  $g(C_r)$  is a function monotonically decreasing with  $C_r$ .

**Corollary 3.2** LSH can find the exact nearest neighbor with a probability at least  $1 - \delta$  with a time complexity  $O(d \log \frac{1}{\delta} n^{g(C_r)} \log n)$  and space complexity  $O(\log \frac{1}{\delta} n^{(1+g(C_r))} + nd)$ .  $l$ , the number of hash tables needed, is  $l = O(\log \frac{1}{\delta} n^{g(C_r)})$ .

Dataset	Dimensionality ( $d$ )	Sparsity ( $s$ )	Relative Contrast ( $C_r$ ) for $p = 1$
SIFT	128	0.89	4.78
Gist	384	1.00	1.83
Color Hist	1382	0.027	3.19



## Relative Contrast and PCA Hashing

### Linear Hashing

$$h(x) = \text{sgn}(w^T x + b), w \in \mathbb{R}^d$$

- Suppose  $b = E[x] = 0$

- Want to find  $w$  such that relative contrast of projections is maximized

Projected distance to nearest neighbor  $(w^T q - w^T x_{\text{nn}})^2$   
Expected distance to a random point  $E_x(w^T q - w^T x)^2$

$$\hat{w} = \arg \max_w \frac{w^T E_q [\sum_i (q - x_i)(q - x_i)^T] w}{w^T E_q [(q - x_{\text{nn}})(q - x_{\text{nn}})^T] w} = \frac{w^T S_x w}{w^T S_{\text{nn}} w}$$

$$\hat{w} = \arg \max_w \frac{w^T S_x w}{w^T S_{\text{nn}} w} \Rightarrow \text{eigenvect}(S_{\text{nn}}^{-1} S_x)$$

If distribution of nearest neighbors is isotropic

$$S_{\text{nn}} \approx \alpha I$$

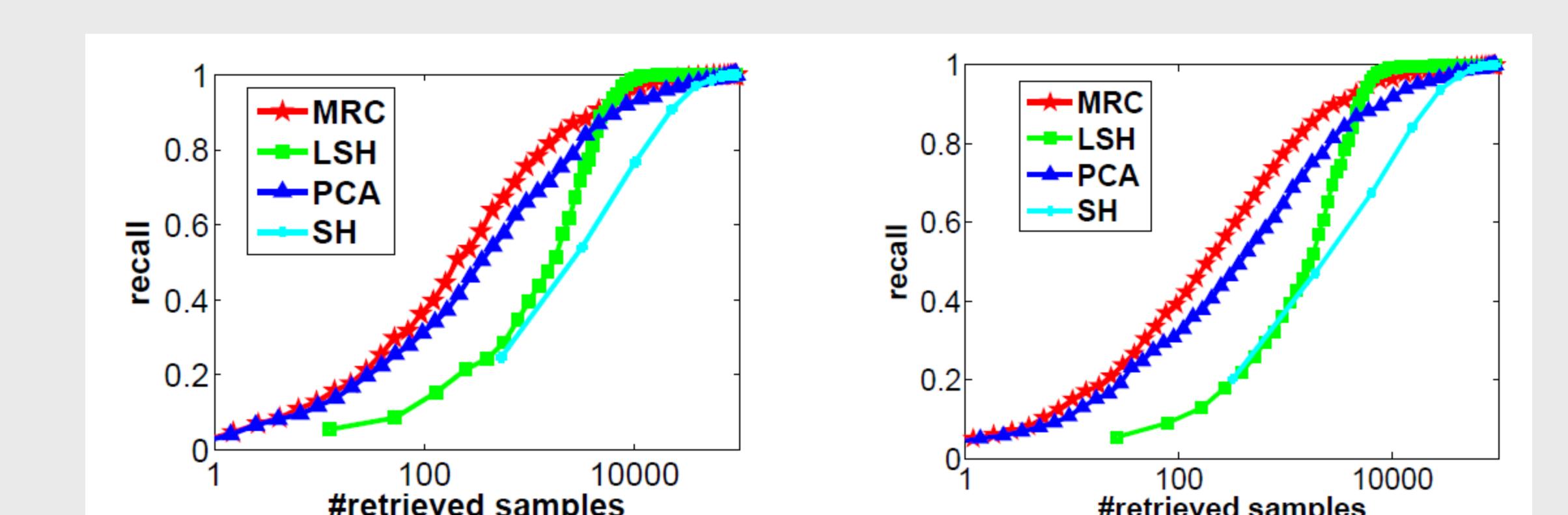
If queries have the same distribution as data,  $q \sim x$

$$S_x \approx \beta \Sigma_x$$

$$\hat{w} = \text{eigenvect}(\Sigma_x)$$

PCA-directions !

Recall of 1-NN with PCA



higher dimension → bad

good ← Lower p

Large database → good